Torwards an Adaptive Middleware for Opportunistic Environments: a Mobile Agent Approach



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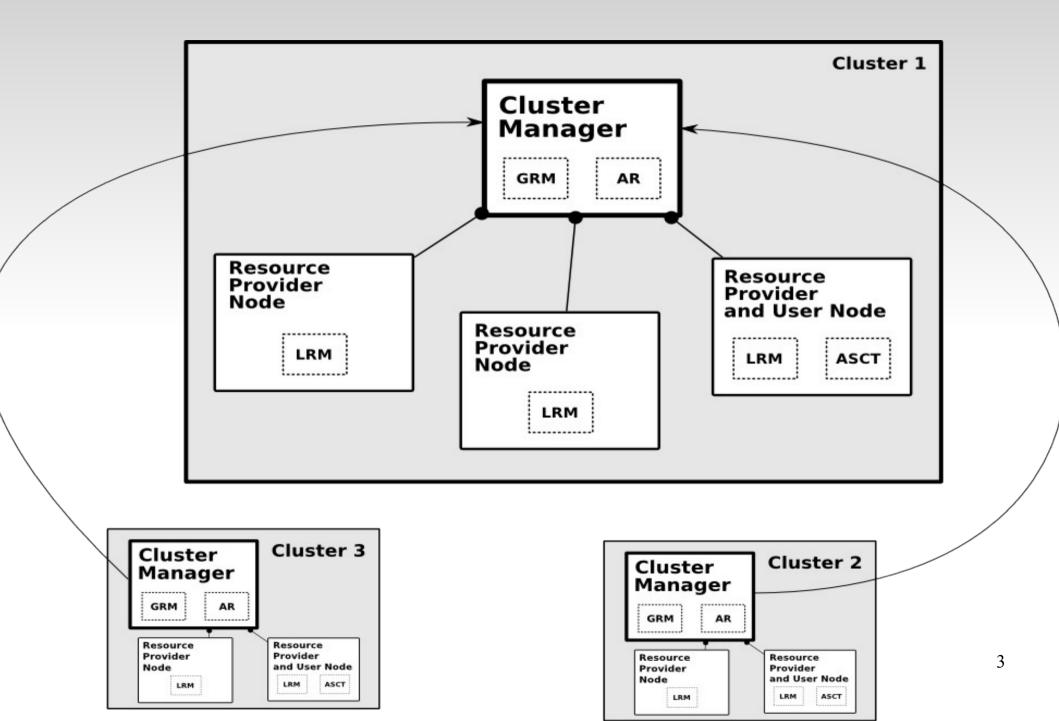




Introduction

- Opportunistic Grids: usage of idle time of nondedicated resources
 - High heterogeneity of resources
 - Failure rate is higher than in dedicated environments
 - Resources "fail" all the time
- InteGrade: Grid middleware for opportunistic grids
 - Usage of idle power from personal computers
 - Architecture: federation of clusters
 - Sequential, parametric, BSP, and MPI applications

Integrade Architecture

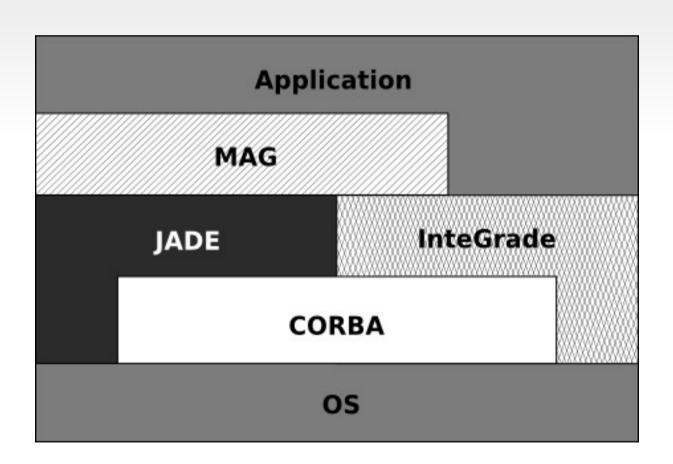


MAG

- MAG: Mobile Agents for Grid Computing
 - Built on top of the InteGrade architecture
 - JADE: agent platform to provide agent communication and life cycle monitoring
 - Mobile agents as a good alternative to build faulttolerance mechanisms
 - Cooperation, autonomy, platform independent, reactivity, and mobility
 - Replication, checkpointing, and retrying for sequential and parametric Java applications

MAG

- MAG: Mobile Agents for Grid Computing
 - Layers of the InteGrade/MAG middleware



Motivation

 Fault-tolerance is essential, specially when executing long-running parallel applications

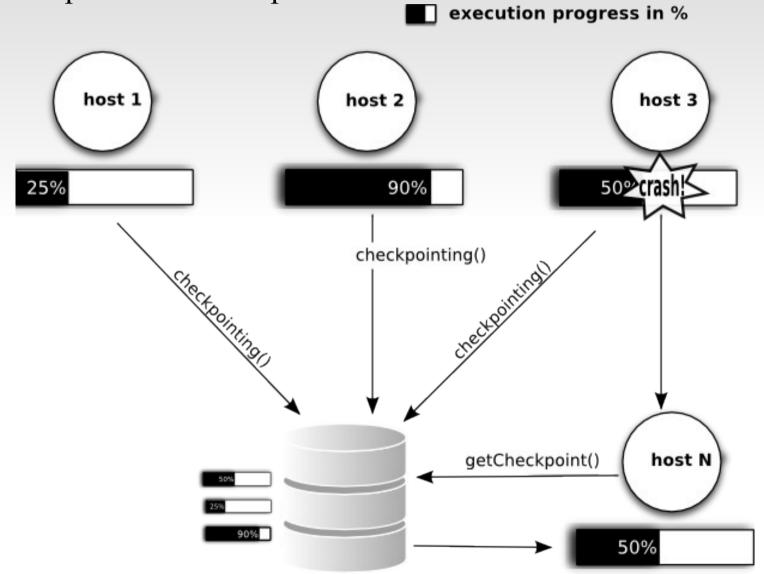
- Failure of a single node require restarting the application from the beggining
- Replication and checkpointing can be used as faulttolerance mechanisms

Fault tolerance on MAG

- MAG supports retrying, replication, and checkpointing of applications
- Weak points
 - These mechanisms operate solely
 - All replicas perform checkpoint periodically
 - If the most advanced replica crashes, its checkpoint will not be reused by other replicas
 - These mechanisms do not perform any automatic adjustments to adapt themselves to changes in resource availability
 - Ex.: nodes leaving and joining the grid

Fault tolerance on MAG

 Recovery: when a replica crashes, it resumes its execution from its last particular checkpoint

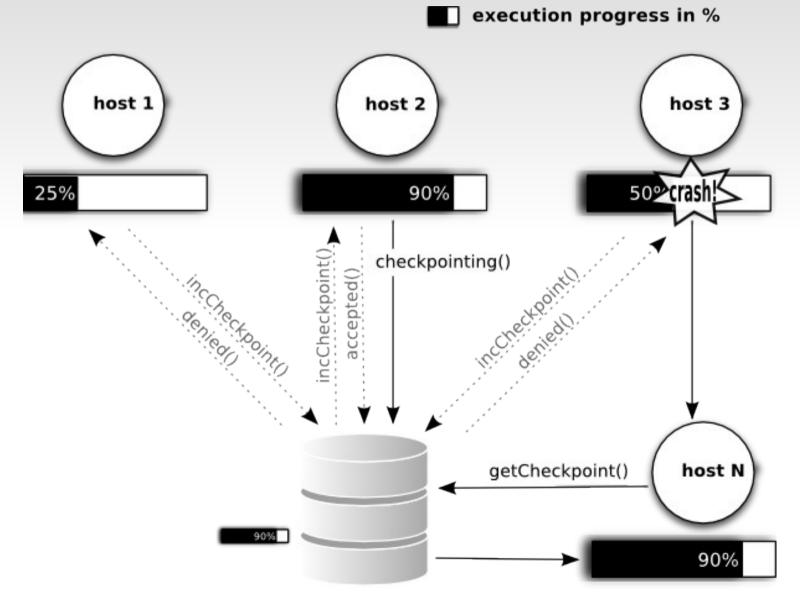


Unified Checkpoint

- Replicas periodically send information about their execution progress and only the most advanced replica is authorized to perform checkpointing
 - The application programmer must manually invoke a superclass method which increases a counter
 - When the replica hits a checkpoint, it sends only the value of the counter
 - The Stable Storage component compares this value to the ones sent by other replicas
 - If this value is the hightest, it sends a message to the replica requesting the checkpoint

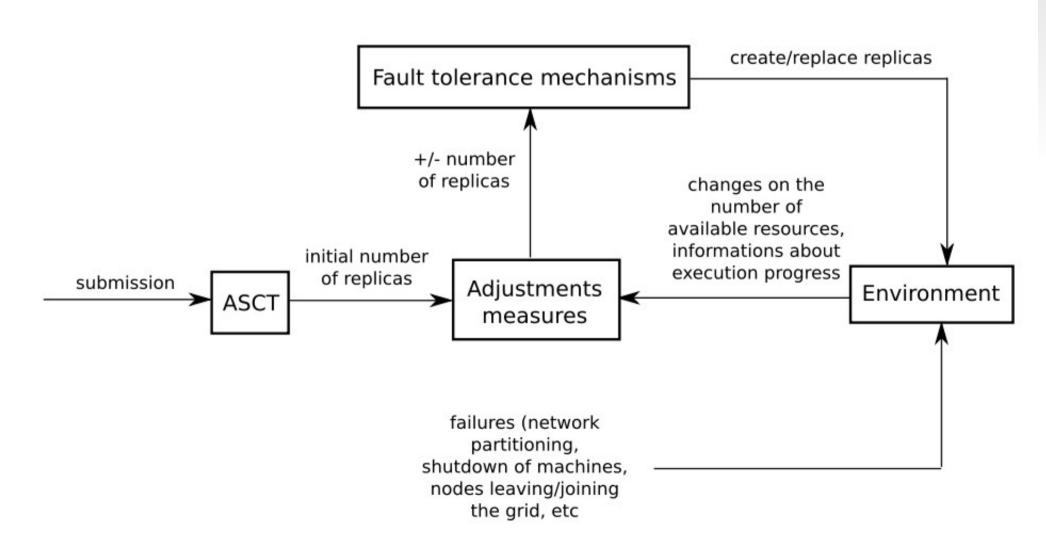
Unified Checkpoint

• Recovery: when a replica crashes, it resume its execution on another machine from the checkpoint of the most advanced replica



Replica Replacement

- Nodes are leaving and joining the grid constantly
- Slow replicas are migrated to improve performance
- Feedback system model



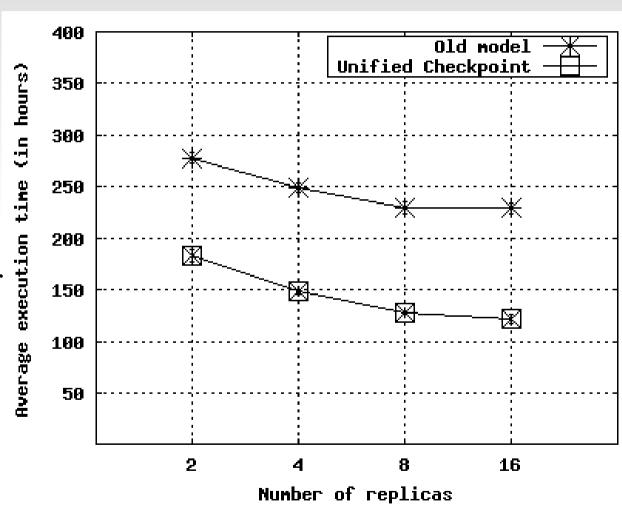
Replica Replacement

- How slow replicas are replaced?
 - StableStorage also checks for slow replicas when comparing replica progression counters
 - If the ratio between a replica counter and the highest counter is below a predefined value, the StableStorage sends a message to the replica requesting its migration to another node
 - After the migration, the replica resumes its execution from the checkpoint of the most advanced replica

- Focus: execution time
- Simulation parameters: failure rate, MTBF (mean time between failures), downtime and number of replicas
- Cluster environment with 100 heterogenous machines connected by a 100Mbps network
- Task model (GridSim Toolkit):
 - 604,8 x 10⁶ MI (millions of instructions)
 - Binary size of 320KB and ouput size file of 15,6KB
 - At least 105 hours of execution

- Simulation scenario built to represent a very inhospitable environment to distributed processing
 - Ex: Student laboratories with machines being regularly turned off and rebooted
 - Fixed 60 minutes as the MTBF
 - 24 failures per day distributed in 100 machines
 - Downtime (average): 30 minutes
- We ran the simulation scenario 40 times with different number of replicas: 2, 4, 8, and 16
 - Compute the average execution time and 95% confidence interval (t-Student distribution)

- Potencial advantage of adopting unified checkpoint happens independently of the number of replicas used
- In all cases: execution times at least 34% lower
- Maximum difference with 16 replicas: 47% lower
- Amount of time saved varies between 95 and 107 hours



Experiments

- Focus: execution time and CPU/Memory consumption
- We submitted a Java application that calculates the approximate value of Pi in an iterative process
 - CPU intensive
 - Could take days of execution (it depends on the input)
 - Many invocations to the checkpoint mechanism
- 16 replicas with all the fault-tolerance mechanisms activated

Experiments

• Execution environment: 17 machines connected by a local Fast Ethernet network (100Mbps)

Machine	Processor	RAM/Swap	OS/Arch	Kernel Version	Distribution
villa	AMD 2.0 GHz	$1~\mathrm{GB}/1.5~\mathrm{GB}$	Linux i686	2.6.22-14-generic	Ubuntu 7.10 (gutsy)
ilhabela	AMD 2.0 GHz	1 GB/1.5 GB	Linux i686	2.6.22.14-generic	Ubuntu 7.10 (gutsy)
taubate	AMD 2.0 GHz	3 GB/768 MB	Linux x86_64	2.6.22.14-generic	Ubuntu 7.10 (gusty)
giga	Intel 3.0 GHz	2 GB/2 GB	Linux i686	2.6.22.14-generic	Debian 5.0 (lenny)
orlandia	AMD 2.0 GHz	$1~\mathrm{GB}/640~\mathrm{MB}$	Linux i686	2.6.22.14-generic	Ubuntu 7.10 (gutsy)
motuca	AMD 2.2 GHz	$1.5~\mathrm{GB/2~GB}$	Linux x86_64	2.6.10	Debian 5.0 (lenny)
mercurio	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
venus	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
terra	AMD 1.4 GHz	$1~\mathrm{GB}/1.5~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
marte	AMD 2.0 GHz	$1~\mathrm{GB}/2~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
jupiter	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
saturno	AMD 1.4 GHz	1 GB/1.2 GB	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
urano	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
netuno	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
plutao	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
hubble	AMD 1.4 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-9-generic	Ubuntu 8.10 (intrepid)
callisto	AMD 1.5 GHz	$1~\mathrm{GB}/0~\mathrm{GB}$	Linux i686	2.6.27-7-generic	Ubuntu 8.10 (intrepid)

- Application execution time
 - without Unified Checkpoint: 63 hours and 30 minutes
 - whit Unified Checkpoint: 40 hours and 42 minutes
- Middleware memory consumption (Jconsole tool)
 - without Unified Checkpoint: 17MB (avg), 30MB (peak)
 - with Unified Checkpoint: 20MB (avg), 34MB (peak)
- Middleware CPU consumption (orlandia machine)
 - with or without Unified Checkpoint: 0,8%

Conclusions and ongoing work

- Unstable and highly heterogeneous environments like opportunistic grids can benefit from dynamic fault-tolerance mechanisms to improve the execution of sequential and parametric applications.
- Checkpointing and replication can work together to reduce resource consumption and improve application execution, and we showed that the Unified Checkpoint is a viable solution.
- Currently, we are investigating other adaptive mechanisms:
 - increase/decrease number of replicas according to failure rate, free resources, and resource competition;
 - changing the checkpointing interval according to failure rate and checkpoint size.

Questions?

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For more information, please visit the project site:

ccsl.ime.usp.br/integrade